

CORADD: Correlation Aware Database Designer for Materialized Views & Indexes

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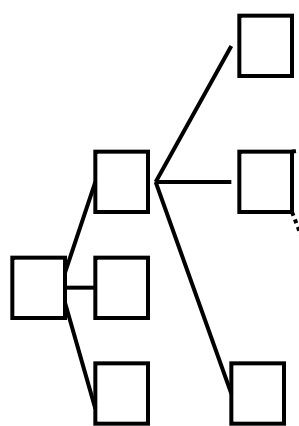
Intro: Data-warehouse (OLAP)

- Very Large Table
Million-Billion tuples
- Analytical Query
Not Very Selective Predicates
- Huge Market
Several \$billion/year

Goal: *Speed Up OLAP Queries*

Problem: Secondary Index

Secondary B+Tree Heap File



<u>state</u>	<u>rowid</u>
MA	
MA	
MN	

<u>state</u>	<u>salary</u>	<u>...</u>
MA	\$25k	
CA	\$20k	
RI	\$30k	
	...	

Bitmap Scan



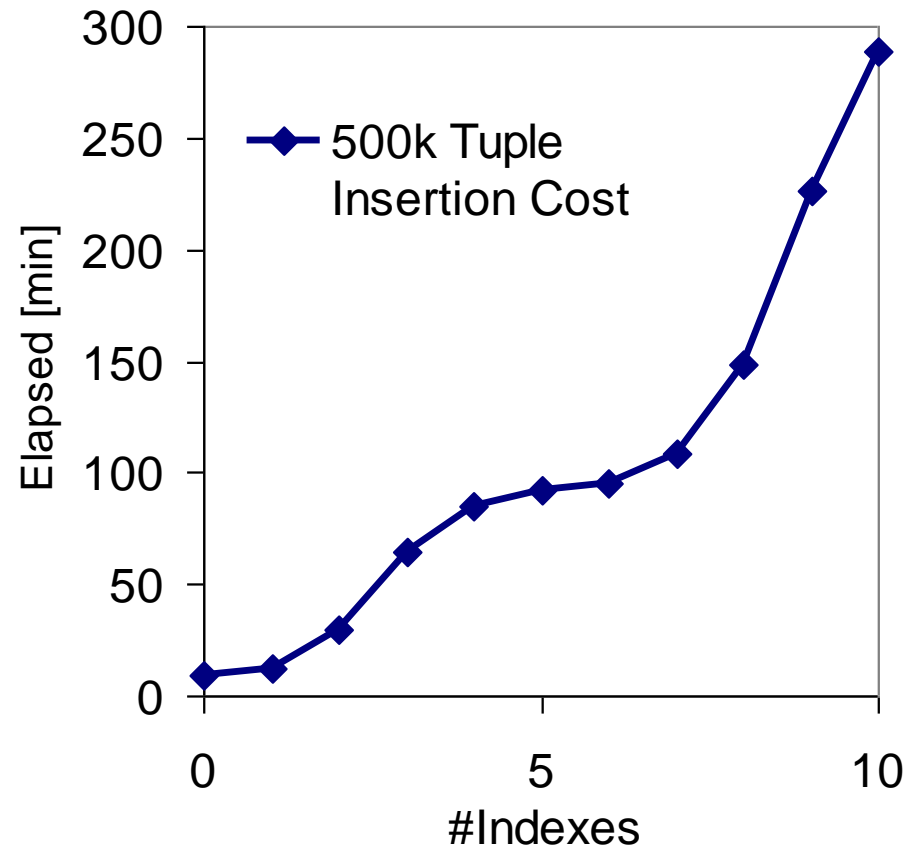
OLAP: Less Selective

≈ Full Table Scan

Secondary Index is Too Slow

Problem: Costs of B+Tree

- Storage Cost
- Maintenance Cost
 - More Pressure on Bufferpool
 - More Eviction from Bufferpool = disk writes



Secondary Index is Too Large

Talk Overview

Idea:

*Exploit Correlations between
Primary and Secondary Indexes*

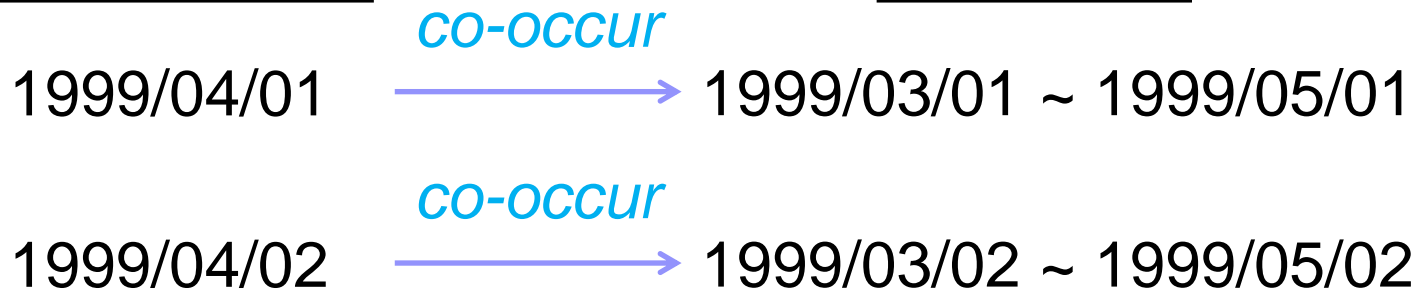
- Correlation and its potentials
- Correlation-Aware Database Designer
- Experiments

What is Correlation?

ORDERKEY	COMMITDATE	SHIPDATE	PRICE
123	1999/04/01	1999/04/06	\$1223
124	1995/09/20	1995/09/19	\$210
125	1997/01/01	1997/01/11	\$6543
...			

COMMITDATE

SHIPDATE



COMMITDATE is correlated with SHIPDATE,
but not with ORDERKEY or PRICE

Correlation and Query Runtime

```
SELECT SUM(PRICE) FROM LINEITEM  
WHERE COMMITDATE = 1995/01/01
```

Secondary
Index

<u>Commitdate</u>
1994/12/31
1995/01/01
1995/01/02
...

Case 1:

Clustered Index

<u>Order Key</u>	Tuple Data
[Greyed out]	

Runtime
(TPCH Scale 20)

150 sec

Case 2:

<u>Ship Date</u>	Tuple Data
[Black highlighted row]	

6 sec

Correlation and Compact Index

Secondary indexes can be orders of magnitude smaller with correlation.

- Correlation Map (CM) [VLDB'09]
Stores soft functional dependencies
- IBM DB2 BHUNT
Stores algebraic relationships b/w attributes
- SQLServer Datetime Correlation Optimization
Stores co-occurring datetime pairs

Correlation Map (CM)

People Table
(Clustered on State)

```
SELECT AVG(Sal.) FROM People  
WHERE City = Boston
```

PID	Sal.	City	<u>State</u>
1	\$25K	Boston	MA
2	\$90K	Cambr.	MA
3	\$80K	Jackson	MN
4	\$30K	Boston	NH
5	\$50K	Manch.	NH
6	\$20K	Toledo	OH

Correlation Map
(City -> State)

<u>City</u>	<u>State</u>
Boston	{MA,NH}
Cambr.	{MA}
Jackson	{MN}
...	...

1,000x Smaller Than Conventional B+Tree

Correlation in DB Design

- **Good News**

Faster and Smaller (Maintainable)
Secondary Index

- **Bad News**

Too Complex for Human DBAs to
Pick Primary/Secondary Indexes

Let's Automate Designing

Problem Statement

Automatic Database Designer

- Input: Queries, Space Budget
- Output: Database Objects
 - Materialized Views (MVs)
 - Clustered Indexes
 - Secondary Indexes
- Constraint: $\text{Size}(\text{Objects}) \leq \text{Budget}$
- Objective: $\min \text{TotalRuntime}(\text{Queries})$

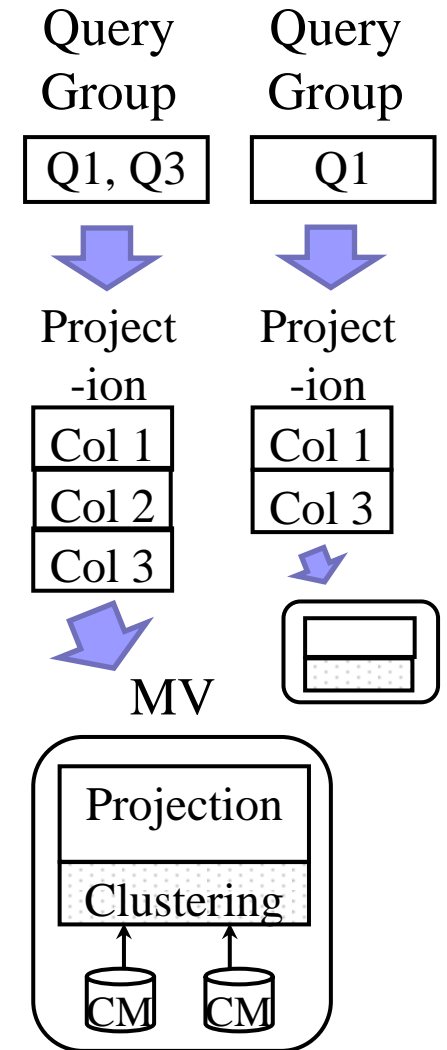
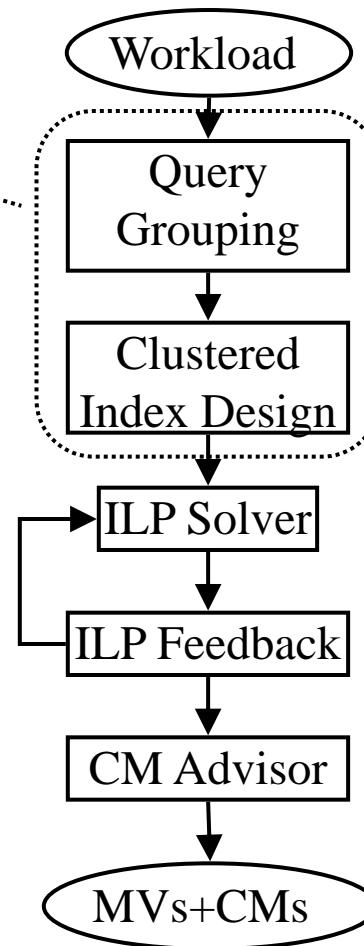
Automatic Database Designer

- Plethora of work In Industry/Academia
- Key Component in Today's DBMS
 - Microsoft SQLServer Database Tuning Advisor
 - IBM DB2 Design Advisor

Still, No Existing Database Designer Exploits or is aware of Correlation

CORADD Framework

- MV Candidate Generator
 - Projection (Columns)
 - Clustered Index
- Correlation Aware Cost Model
- ILP/ILP Feedback
- CM Advisor



MV Candidate Generator

- MV = Projection + Clustered Index
- Projection Design
 - Choose Queries To Serve
 - Maximize Similarity Of Queries (Predicates)
- Clustered Index Design
 - Choose Clustered Index On Projection
 - Maximize Correlation with Secondary Indexes (CMs)

Projection: Query Grouping

■ Selectivity Vector

<u>Query</u>	<u>year</u>	<u>yearmonth</u>
Q1	0.15	→ 0.15
Q2	← 0.15	0.013

Predicates
Correlated
year = '95
yearmonth = '94Dec

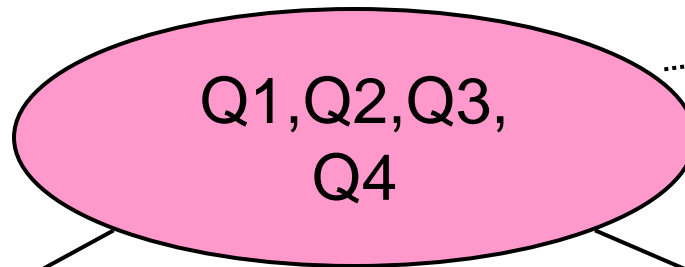
- *Selectivity Propagation*

■ k-means to Group Similar Queries

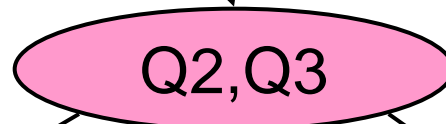
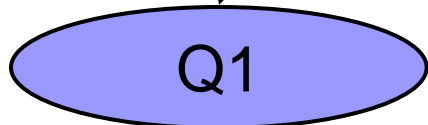
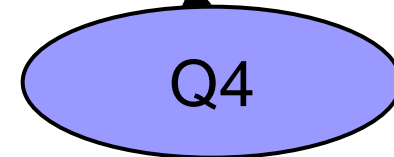
$$\text{dist}(\vec{v}_1, \vec{v}_2) \equiv \sqrt{\sum_{attr} (v_1[attr] - v_2[attr])^2}$$

Clustered Index Design: Split

Split by
k-means (k=2)

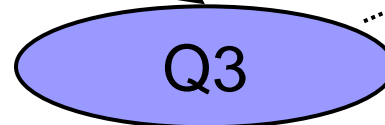
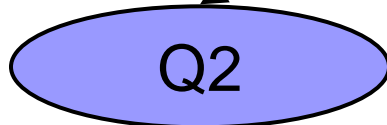


Projection

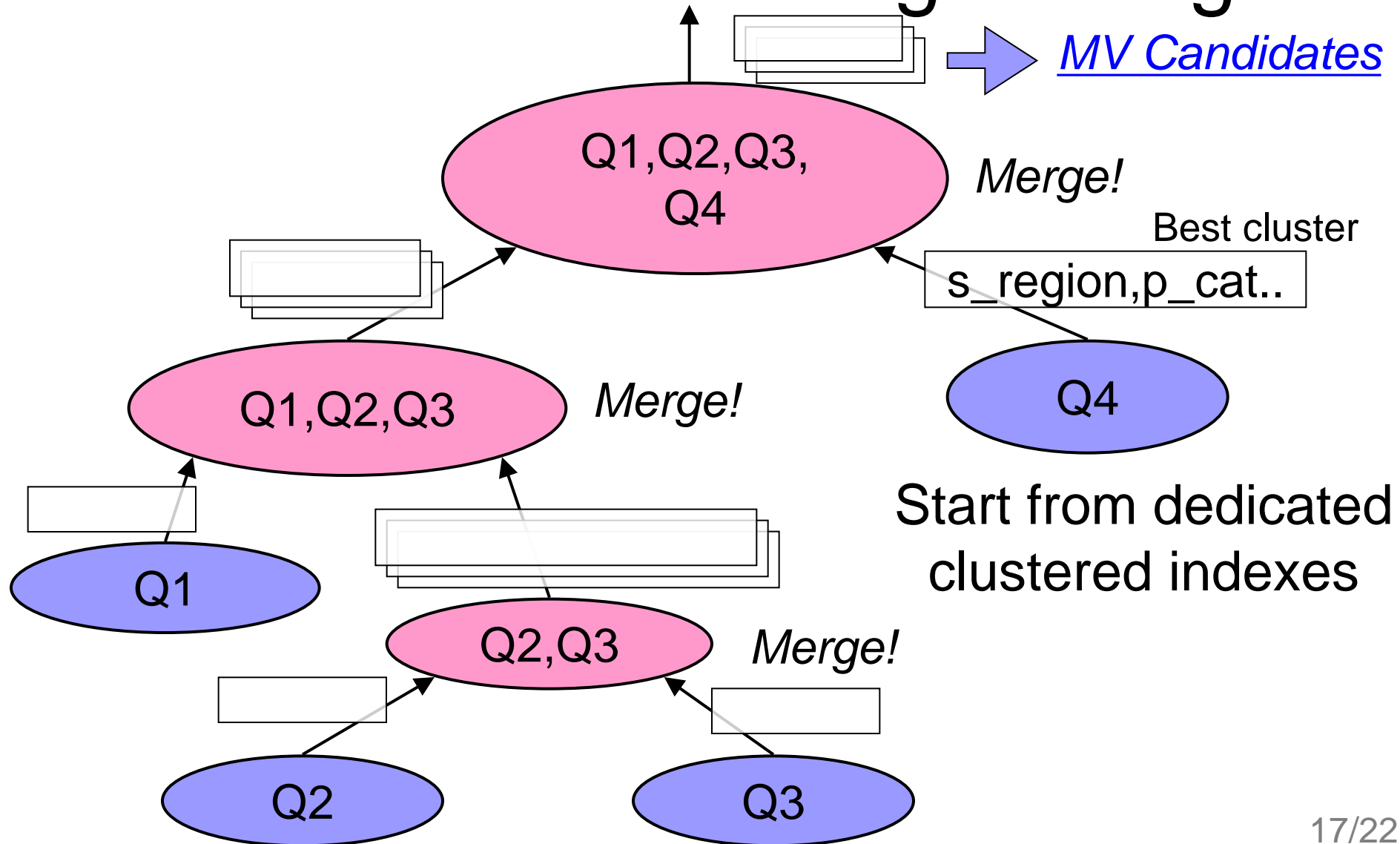


Dedicated MV

Easy to design
best clustered index



Clustered Index Design: Merge



Experiments

■ Environment

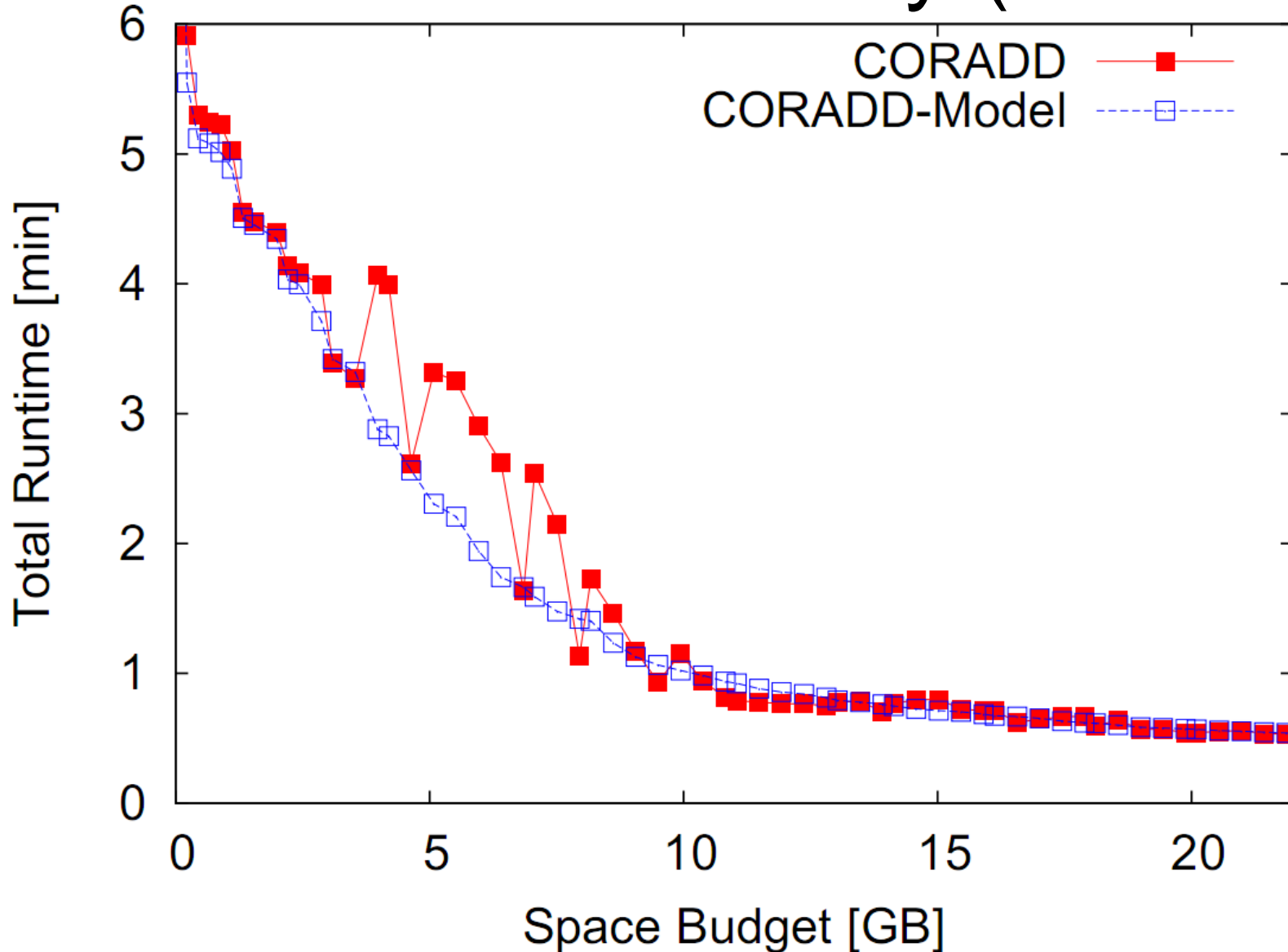
- DBMS X on Windows Server Enterprise Ed.
- Quad-core, 4GB RAM, 10k rpm HDD

■ Dataset/Workload

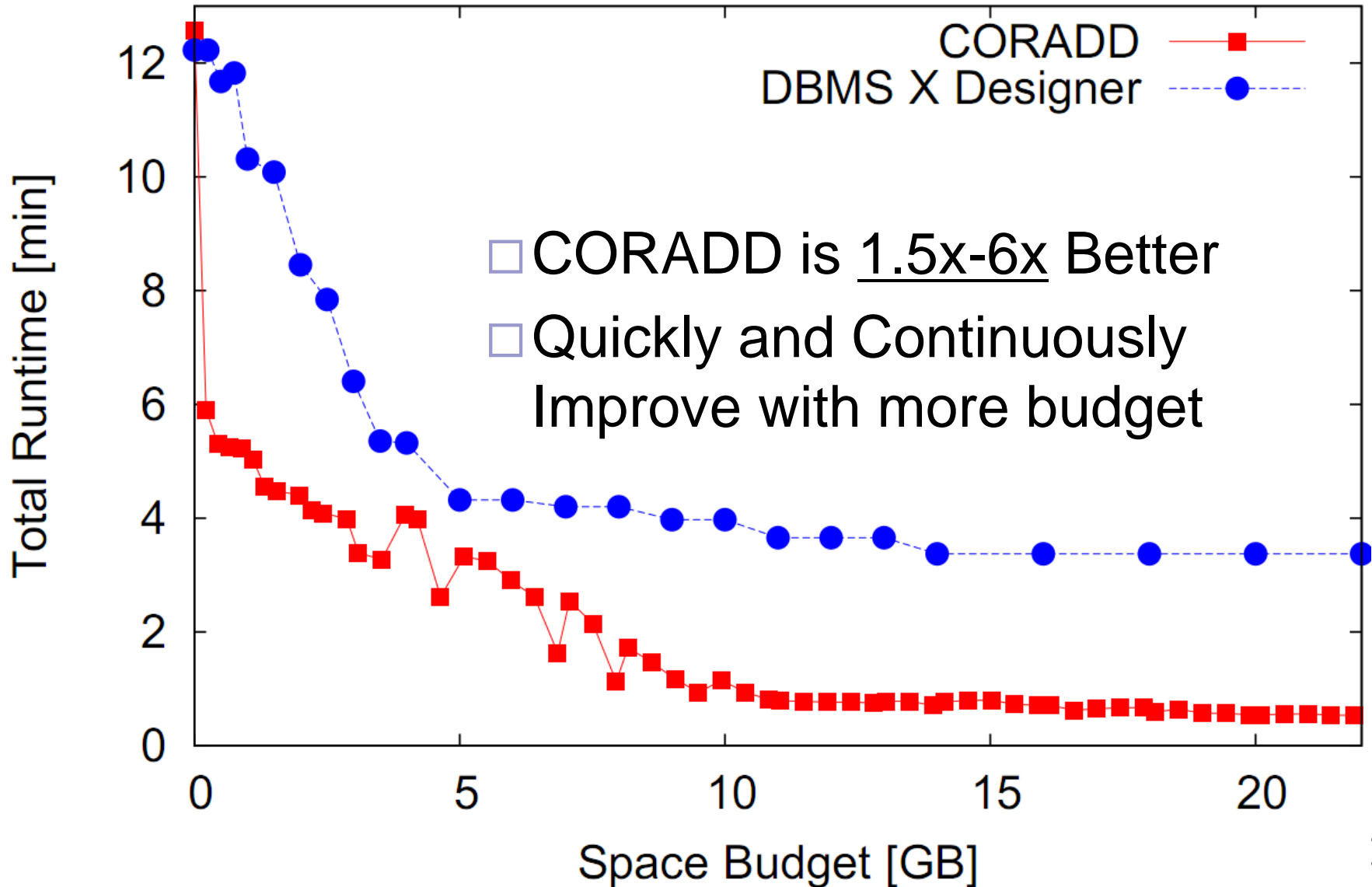
- OLAP Council APB-1 (45M tuples)
- Star Schema Benchmark (24M tuples)

■ Compared w/ DBMS X's own DB Designer

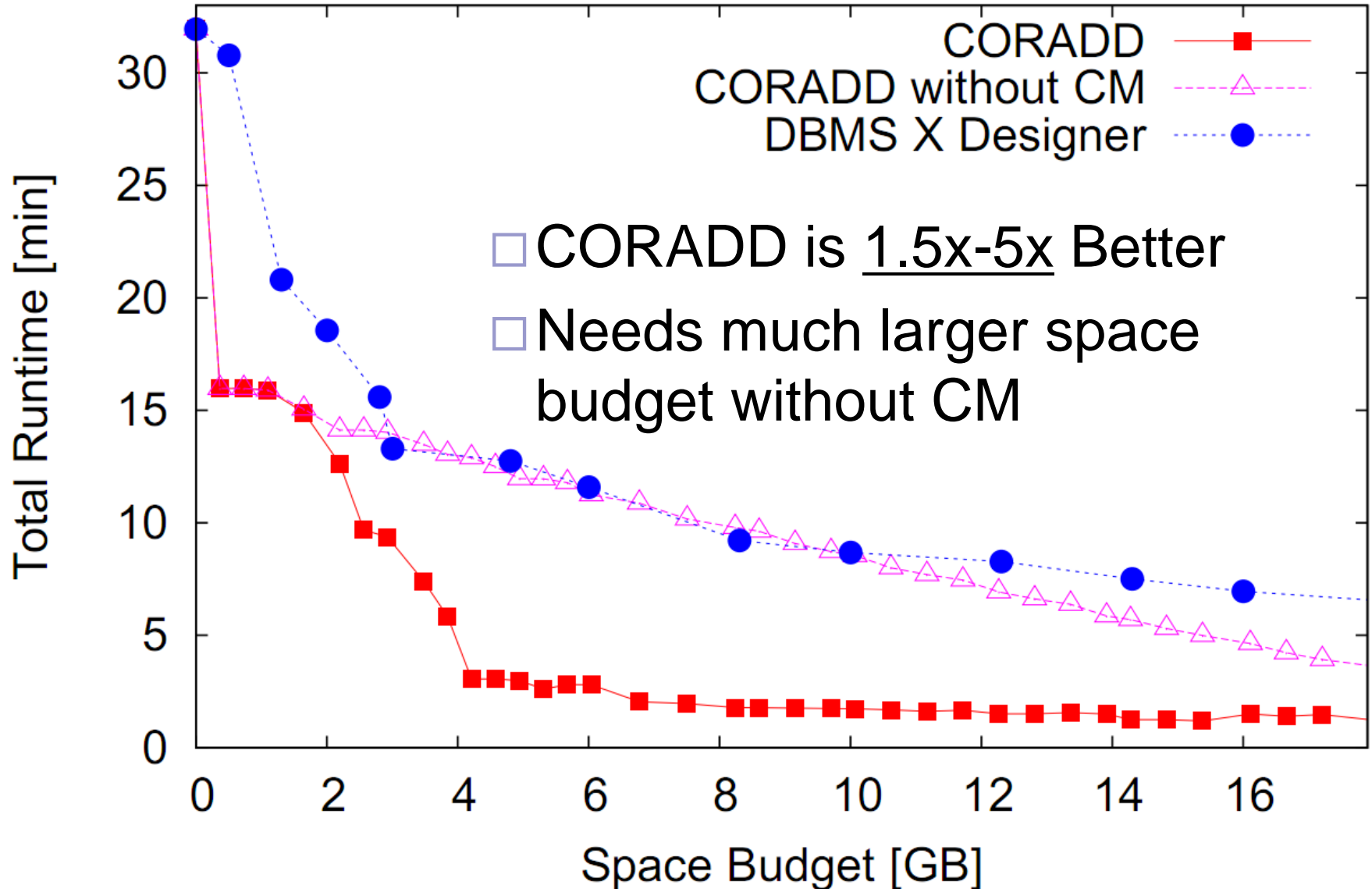
Cost Model Accuracy (APB-1)



Real Runtimes on APB-1



Real Runtimes on SSB



Conclusion

- Index Correlation
 - Correlation = Faster & Smaller Index
 - Choice of Clustered Index matters!
- Correlation-Aware Database Designer
 - Design MVs based on Selectivity Vector
 - Maximize Correlation between Clustered/Secondary Index
 - 5-6x Better Quality to State-of-the-art